



J-NVM: Off-Heap Persistent Objects in Java

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Non-volatile main memory

new persistent medium (in-between **SSD** and **DRAM**)

Durable

resists reboots, power loss

High-density

smallest DIMM = 128 GB

Byte addressable

persistent memory abstraction

High-performance

low latency (seq. read/write ~ 160/90ns)
high bandwidth (up to 8.10GB/s, *2nd gen*)



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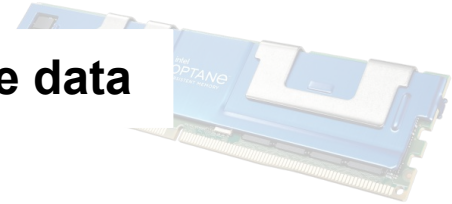
Direct byte-addressability of durable data

Byte addressable

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1- Dramatic throughput and latency improvement for persistent data applications

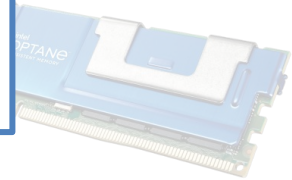
Byte addressability

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High-density

smallest DIMM

1- Dramatic throughput and latency improvement for persistent data applications

Byte addressable

persistent memory abstraction

2- Simpler code bases with single data representation and no file I/Os

High-performance

low latency (seq. read/write ~ 100/90ns)
high bandwidth (up to 8.10GB/s, 2nd gen)



Why Java?

Many data stores & processing frameworks

- Spark, Hadoop, Kafka, Flink, Cassandra, HBase, Elasticsearch, etc.

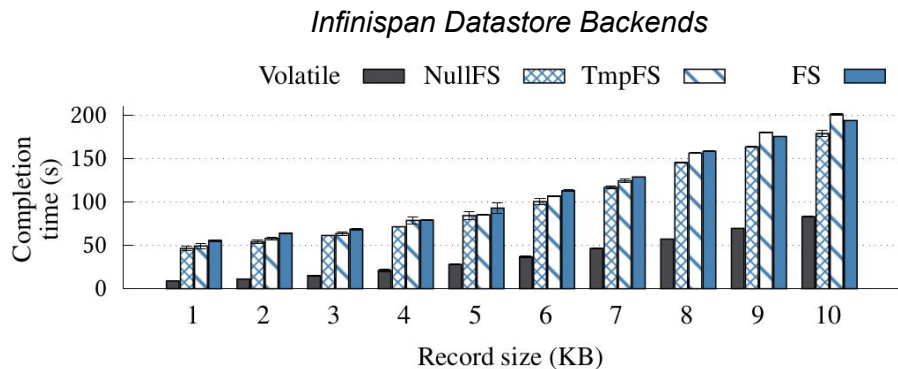
Lack of *efficient* interfaces

- **FS/ext4-dax**

- almost as slow as tmpfs
- dual representation (consistency)
- cost of marshalling

- **PCJ (JNI+PMDK)**

- slower than FS on YCSB benchmark



Varying record size in YCSB-F.

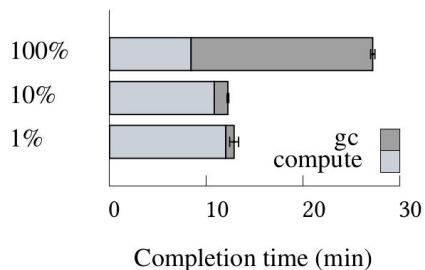
Problematic: Java-native NVMM interface

Prior works: *internal design*

= [Espresso, AutoPersist, go-pmem]

Managed Persistent Objects

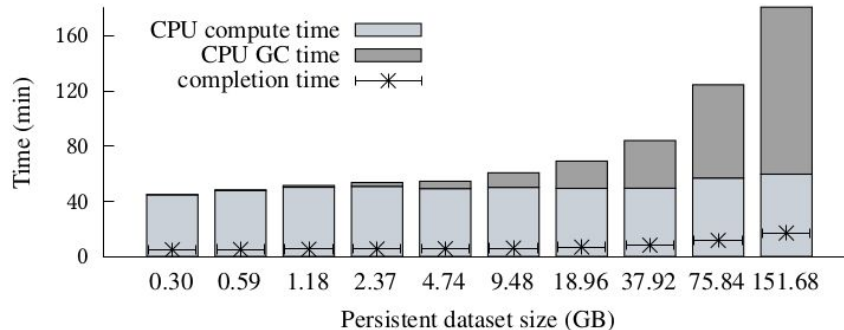
Infinispan datastore in-memory cache ratio



Fixed dataset size - 80GB on heap for 100% cache

Varying cache ratio (YCSB-F)

go-redis-pmem with increasing dataset size



Increasing dataset (YCSB-F, go-pmem)

Features

managed persistent objects

orthogonal persistence
(pnew, @persistentRoot)

heavily-modified runtime

failure-atomic blocks

non-scalable

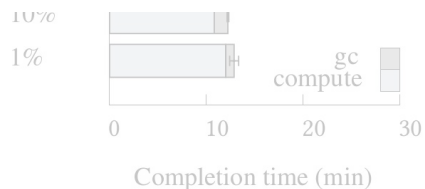
Garbage collectors can not scale to large persistent datasets

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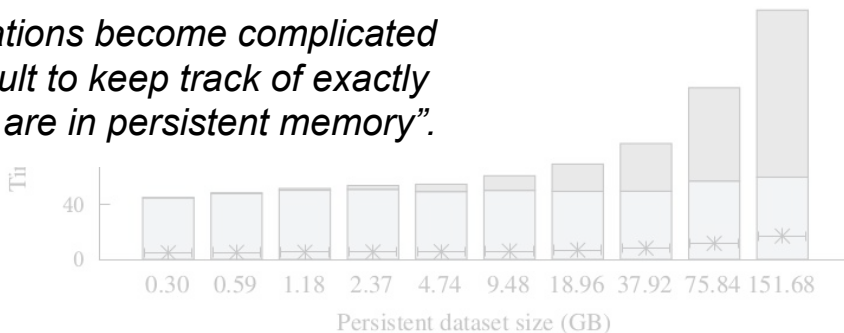
In [go-pmem]: “as the applications become complicated it becomes increasingly difficult to keep track of exactly which variables and pointers are in persistent memory”.



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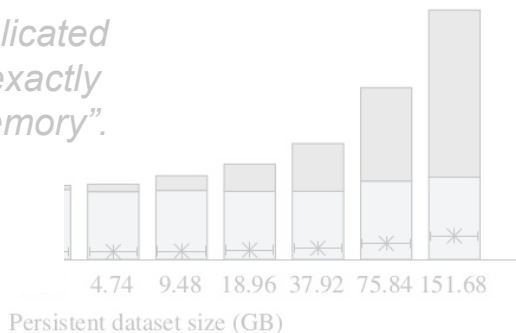
code instrumentation = made whole JVM 51% slower in [Autopersist]

Completion time (min)

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Increasing dataset (YCSB-F, go-pmem)

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(**pnew**, **@persistentRoot**)

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Garbage collectors can not scale to large persistent datasets

Overview

J-NVM = Off-Heap Persistent Objects

Challenges

single data representation

programming model

direct access to NVMM

durability abstraction

scalability (large persistent dataset)



Features

off-heap persistent objects

class-centric model
(code generator + PDT library)

sun.misc.Unsafe

failure-atomic blocks + fine-grained

see evaluation

Outline

Introduction

- NVMM
- why Java?
- prior works
- overview

System Design

- programming model
 - persistent objects
 - code generator
- JPFA
- JPDT

Evaluation

- YCSB benchmark
- recovery

Conclusion

Overview

J-NVM = Off-Heap Persistent Objects

Key idea

each persistent object is decoupled into

- *a persistent data structure*: unmanaged, allocated off-heap (NVMM)
- *a proxy*: managed, allocated on-heap (DRAM)

Programming model - *persistent objects*

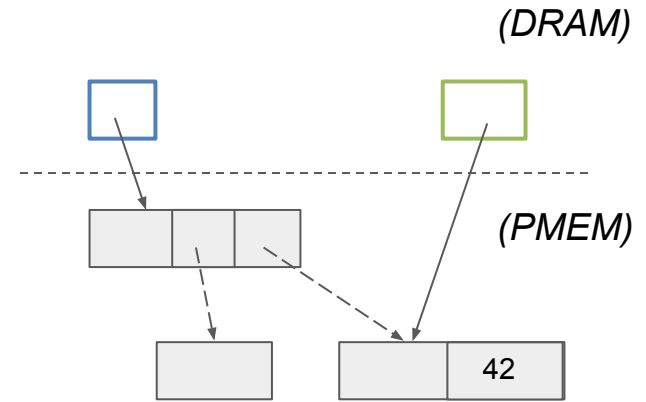
Persistent object is

- a persistent data structure
 - holds object fields
- a proxy
 - holds object methods
 - implement PObject interface
 - intermediate access to pers. data structure
 - instantiated lazily (low GC pressure)

Alive when reachable (from persistent root)

Class-centric model

- safe references thanks to the type system



```
Map root = JNVM.root();
Simple s = root.get("Simple");
s.setX(42);
```

Programming model - *life cycle*

Constructor

- allocate NVMM
- attach persistent data structure

Re-Constructor

- re-attach proxy
- re-build soft state via resurrect()

Destructor

- explicit **JNVM**.free() to reclaim NVMM
- detach proxy
- ready to be GCed

```
Simple s = new Simple(42);
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(DRAM)

(PMEM)

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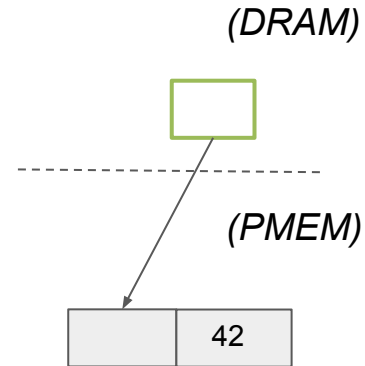
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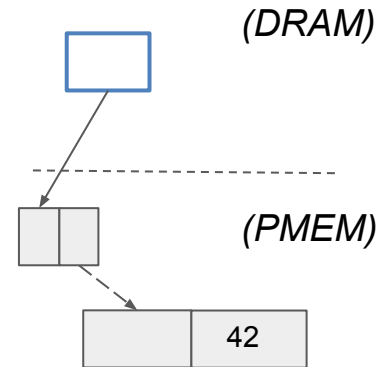
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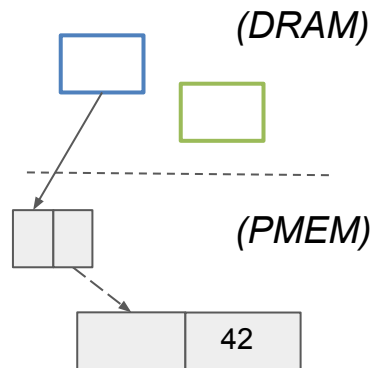
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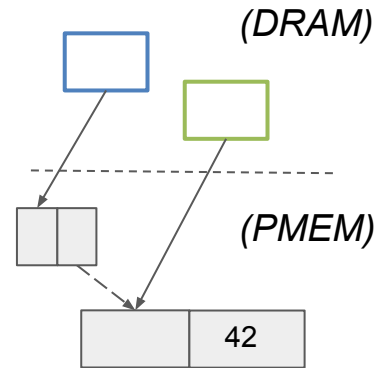
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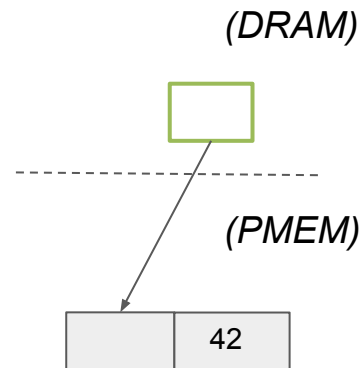
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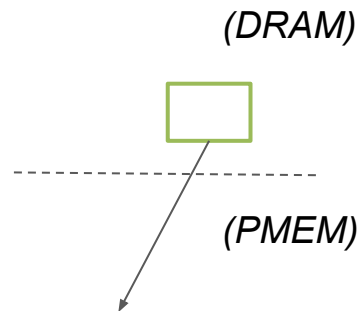
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Overview

J-NVM = Off-Heap Persistent Objects

Tooling

- built-in off-heap memory management for NVMM
- **code generator**: automatic decoupling for POJOs
- **J-PFA**: automatic failure-atomic code
- **J-PDT**: data types + collections for persistent memory
- low-level API (for experts)
- recovery-time GC

Programming model - *code generator*

```
@Persistent(fa="non-private")
class Simple {
    PString msg;
    int x;
    transient int y;

    Simple(int x) {
        this.x = x;
        this.msg = new PString("Hello, NVMM!");
    }

    void inc() { x++; }
}
```

Goals

- class-wide off-heap layout
- generate constructor, re-constructor
- replace (non-transient) field accesses
- wrap non-private methods

```
// transformed code
class Simple implements PObject {
    transient int y;
    long addr; // persistent data structure

    Simple(int x) {
        JNVM.faStart();
        this.addr = JNVM.alloc(getClass(), size());
        setX(x);
        setMsg(new PString("Hello, NVMM!"));
        JNVM.faEnd();
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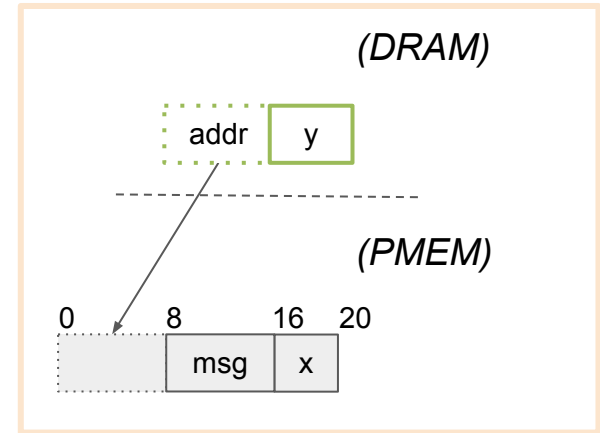
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```
// transformed code (continued)
long addr; // the persistent data structure
long size() { return 12; }
```

```
PString getMsg() { return (PString)
    JVM.readPObject(addr, 0); }
```

```
void setMsg(PString v) {
    JVM.writePObject(addr, 0, v); }
```

```
int getX() {return JVM.readInt(addr, 8);} }
```

```
void setX(int v) {JVM.writeInt(addr, 8, v);} }
```

Goals

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- generate constructor, re-constructor
- replace (non-transient) field accesses
- wrap non-private methods
- generate or transform field accessors

J-PFA

Automatic crash-consistent update

usage = **JNVM**.faStart() *some code* **JNVM**.faEnd()

Per-thread persistent redo-log (inspired by Romulus)

Log new, free and updates

granularity = a block of PMEM

Do *not* log updates to “new” persistent objects
(e.g. allocated within the FA-block)

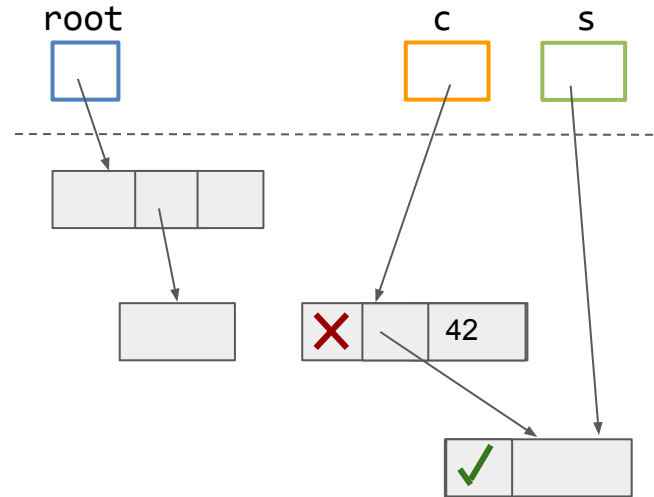
J-PDT + Low-level interface

J-PDT

- drop-in replacement for (part of) JDK e.g., string, native array, map.

Low-level interface

- `unsafe.{pwb,pfence,psync}`
- NVMM block allocator
- recovery time GC (à la Makalu)
- validation = 1 bit in object header
 - makes atomic reclamation easier
 - allows deferring object liveness
 - interpreted on recovery to reclaim reachable invalid objects



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Complex c = new Complex(s);  
root.put("Complex", c);  
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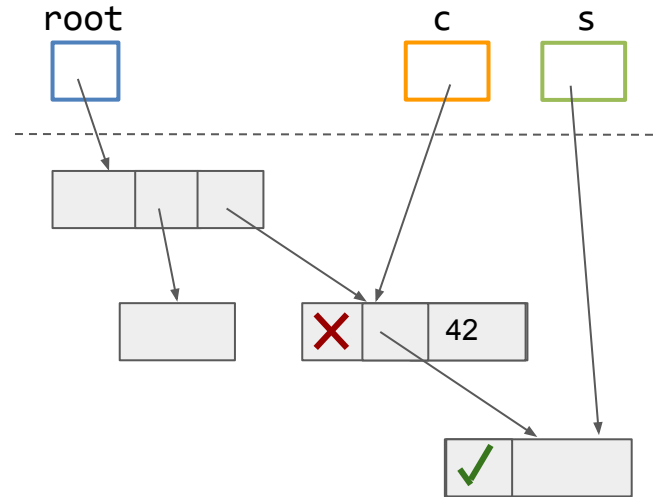
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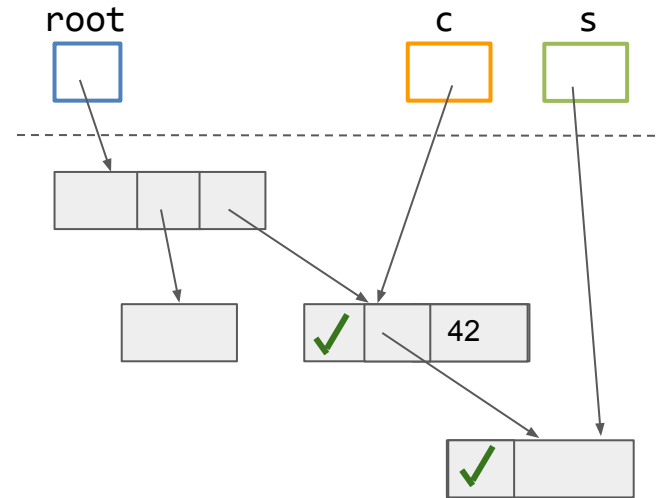
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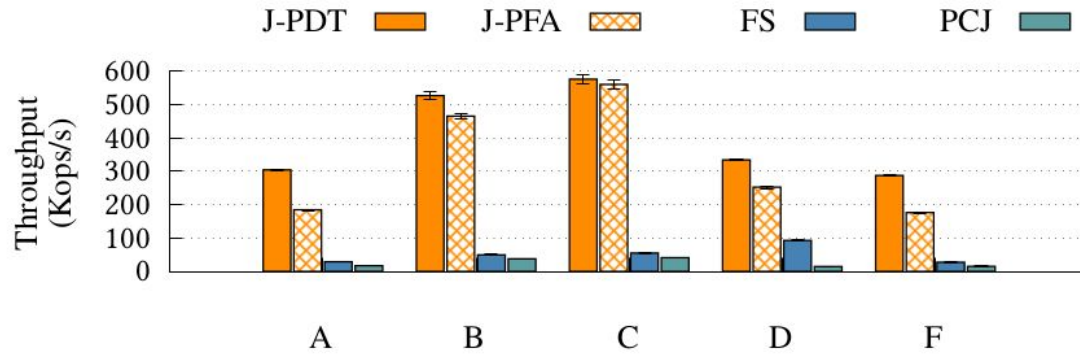
- programming model
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- JPFA
- JPDT

Evaluation

- YCSB benchmark
- recovery

Conclusion

YCSB Benchmark



Durable backends for Infinispan:

- PCJ = HashMap from Persistent Collections Java (JNI + PMDK)
- FS: ext4-dax

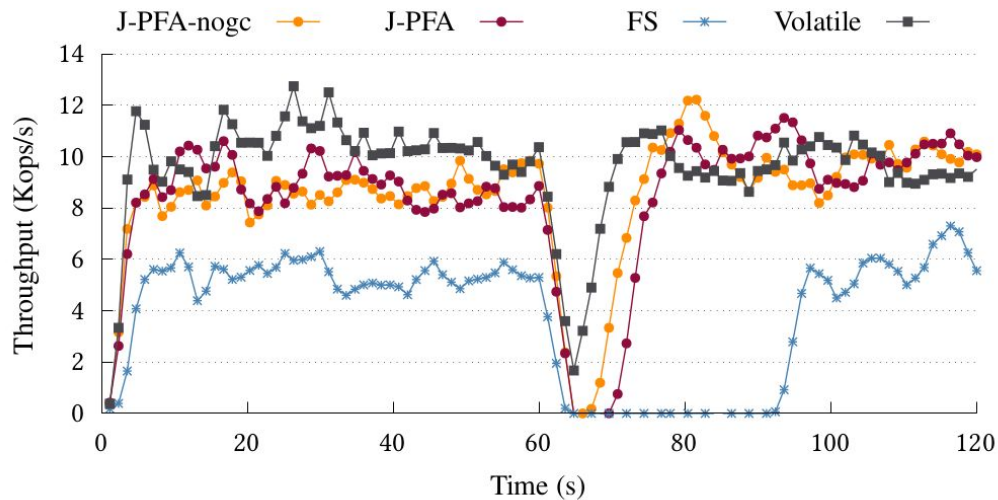
Hardware used:

4 Intel CLX 6230 HT 80-core
128GB DDR4,
4x128GB Optane (gen1)

Takeaways:

- J-NVM up to 10.5x (resp. 22.7x) than FS (resp. PCJ)
- no need for volatile cache

Recovery



TPC-B like benchmark
10M accounts (140 B each)
client-server setting
SIGKILL after 1 min

Takeaways:

- J-NVM is more than 5x faster to recover than FS
- no-need for graph traversal in some cases (e.g., only FA blocks)

Conclusion

J-NVM = off-heap persistent objects

Each persistent object is composed of

- *a persistent data structure*: unmanaged, allocated off-heap (NVMM)
- *a proxy*: managed, allocated on-heap (DRAM)

Pros:

- unique data representation (no data marshalling)
- recovery-time GC (not at runtime, does not scale)
- consistently faster than external designs (JNI, FS)

Cons:

- explicit free but common for durable data
- limited code re-use but safer programming model