



Error-Correcting WOM Codes for Worst-Case and Random Errors

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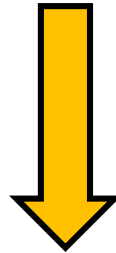
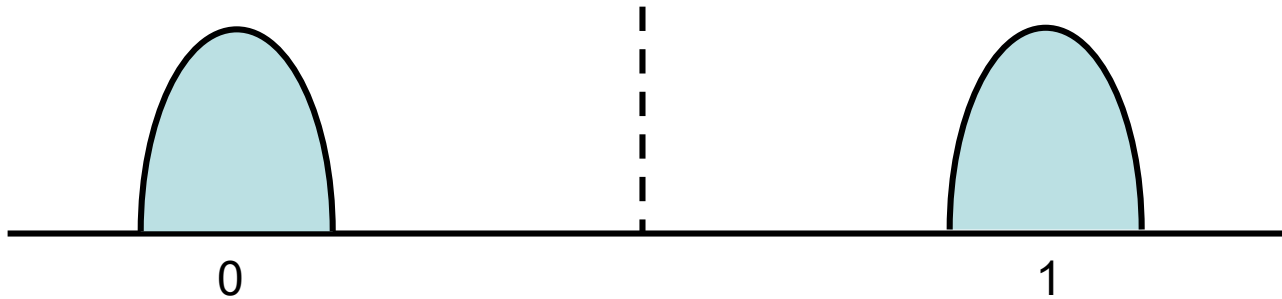
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Multi Level NVM

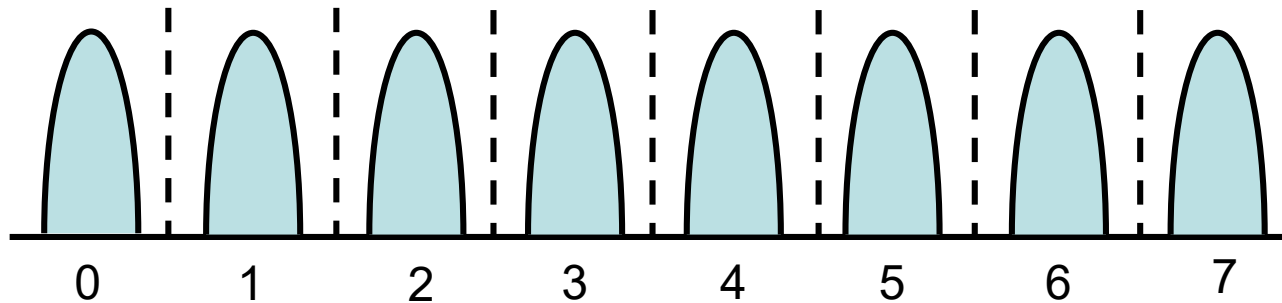


2-Level
Cell



Multi-Level
Cell

Levels:
 $q = 8$



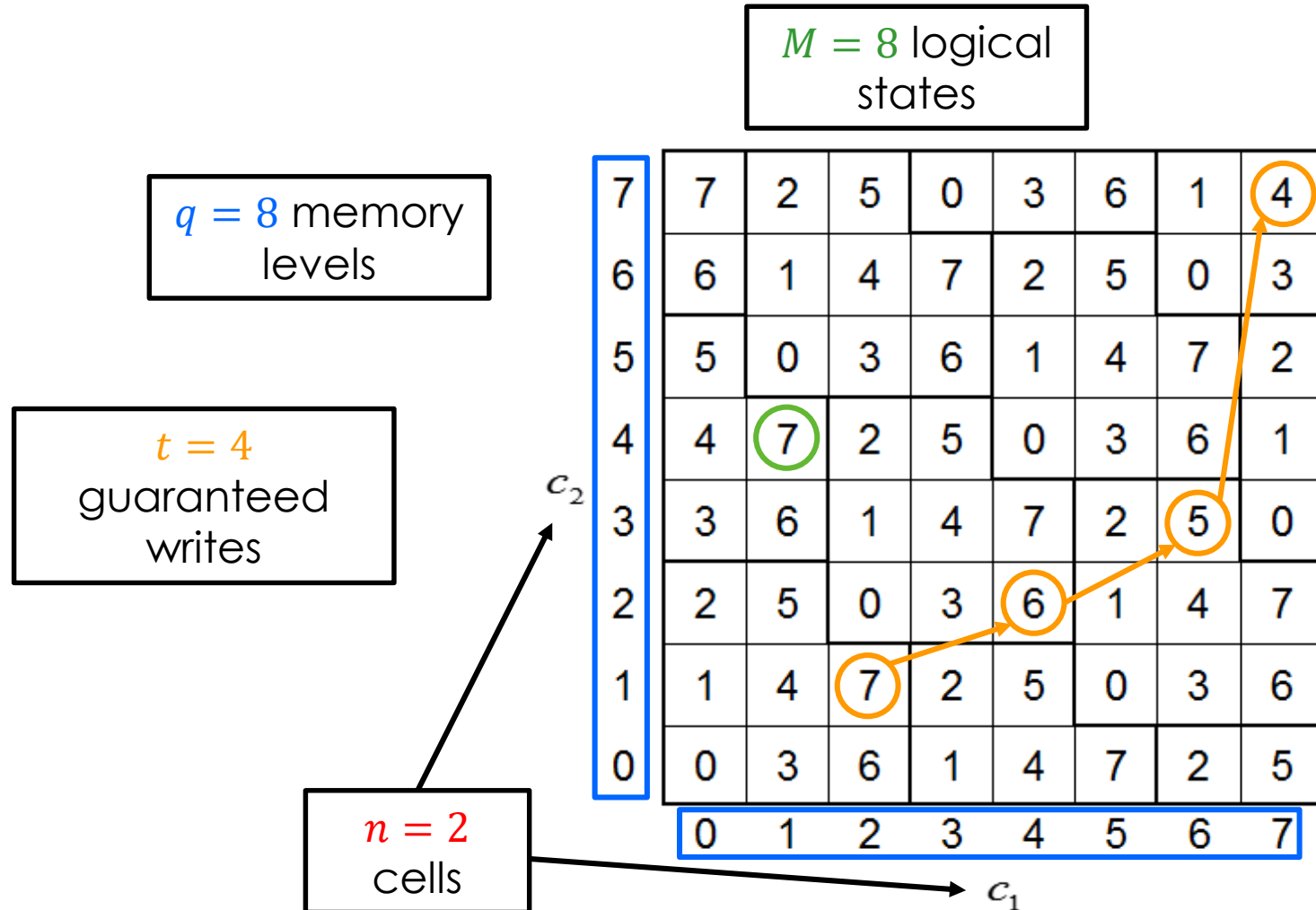


- **The main problem of Flash storage:**
 - Cannot update data in place.
 - Major penalty in access speed and cell wear
- **Reason: cell levels cannot be decreased without erasing a large data unit.**
- **WOM codes: Mitigate this problem [Rivest, Shamir '82]**
 - Allow updating logical data multiple times, without decreasing physical cell levels
- **Potential application: special SSD “hot zone”**
 - File-system journals
 - FTL bypass, e.g. encryption keys, modem information

WOM code - Building Block



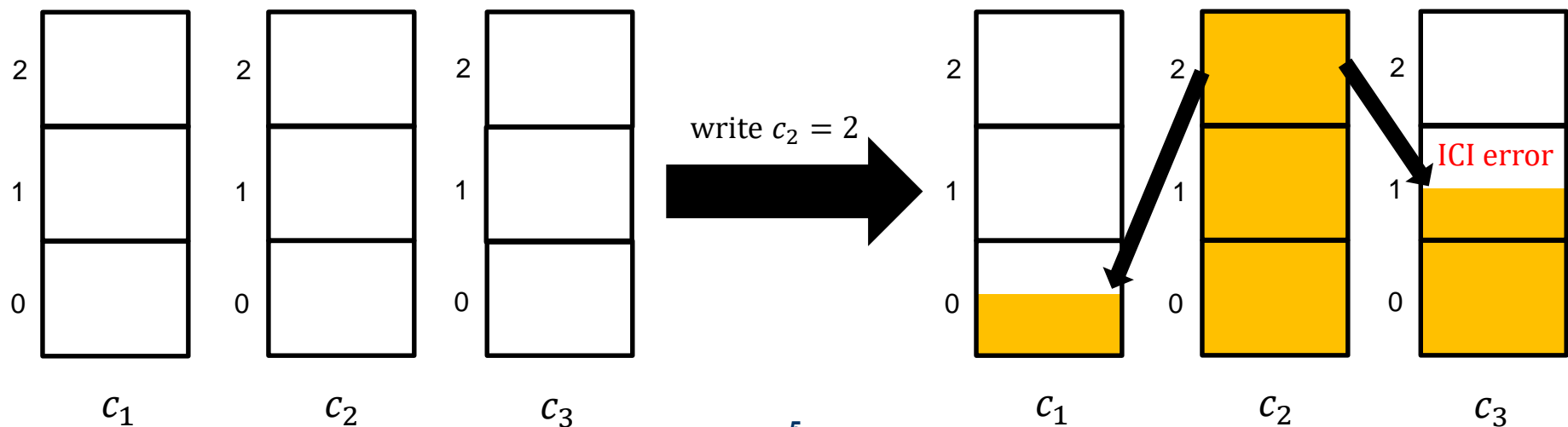
- Example: ($n = 2$, $q = 8$, $t = 4$, $M = 8$)



Problem: Errors and WOM codes



- Memory cells suffer read/write errors
- Re-writing further degrades reliability due to inter-cell interference (ICI)
- Combining WOM codes with error-correcting (EC) codes (ECC) is non-trivial
 - we only have weak or non-practical ECC-WOM codes





■ Design practical error-correcting WOM codes

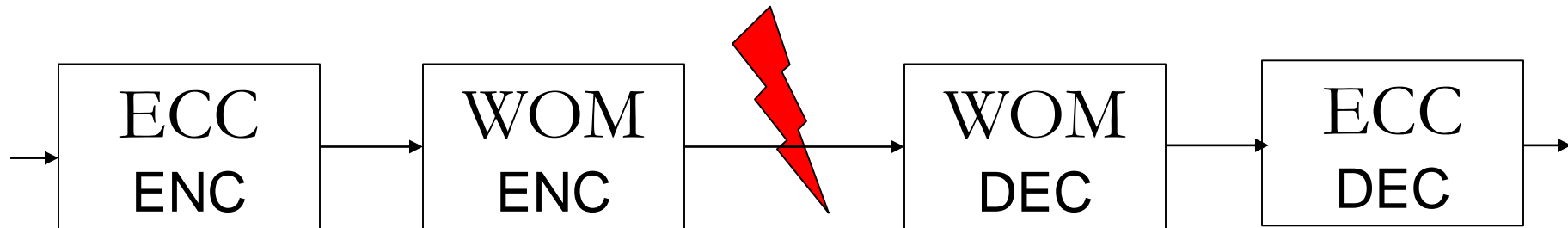
- Low redundancy
- For common memory errors
 1. Guaranteed error correction
 2. Random errors
- Any arbitrary number of errors τ

■ Prior work

- [Yaakobi, Siegl, Vardy, Wolf '12], [Zemor, Cohen '91] – specific τ
- [Jiang, En-Gad, Langberg, Bruck '13] – Polar codes, asymptotic block length

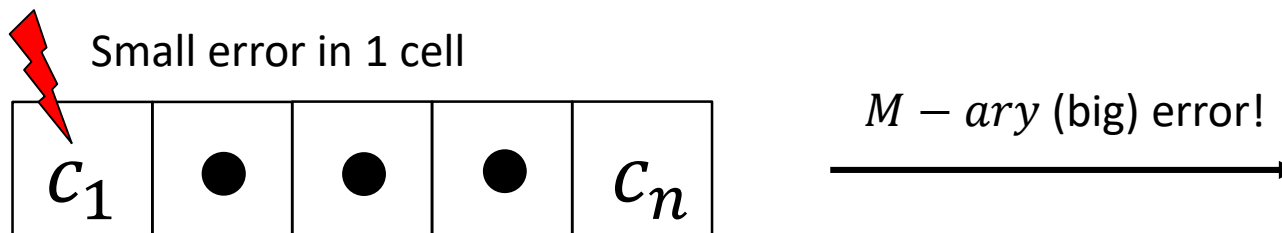


■ Forward concatenation: error propagation

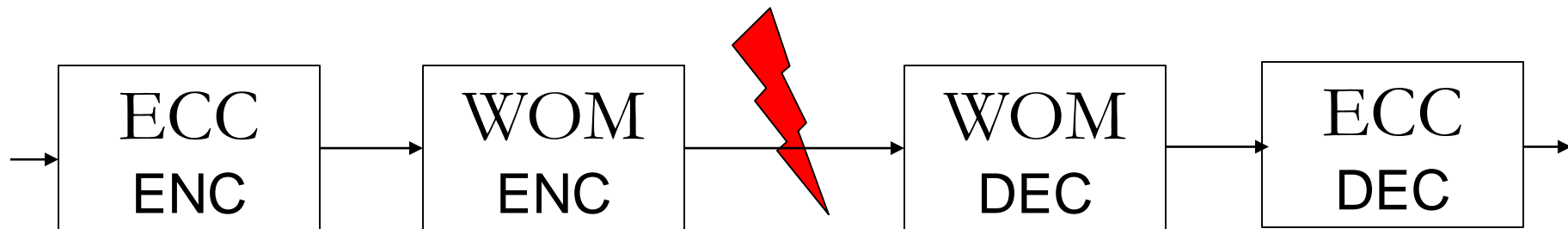


■ Example: (n, q, t, M) WOM

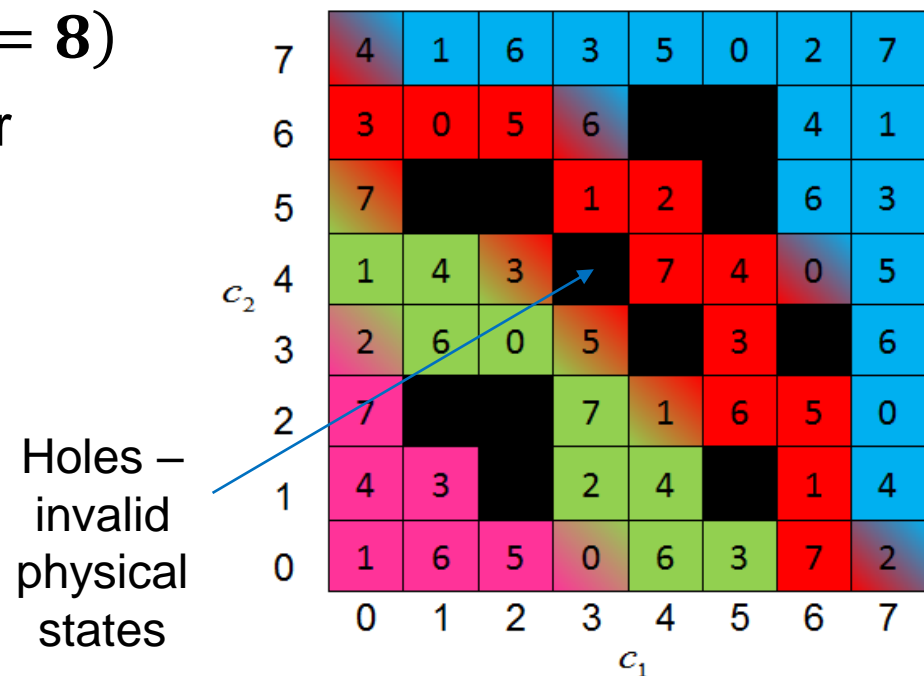
- Small errors in cell levels result in M – ary errors
- WOM is part of the channel



Concatenate ECC with special WOM codes



- ECC “sees” the WOM code as part of the channel
- Example of our special WOM codes
 - $(n = 2, q = 8, t = 4, M = 8)$
 - $k \in \{1, 2, 3, 4\}$ write number
- **Correction Guarantee!**



EC-WOM construction for mag-1 errors



- Mag-1 error results in 4 – *ary* error or erasure
- First correct 4 – *ary* errors

7	c0	a1	d0	b1	c1	a0	b0	d1
6	b1	a0	c1	d0			c0	a1
5	d1			a1	b0		d0	b1
4	a1	c0	b1		d1	c0	a0	c1
3	b0	d0	a0	c1		b1		d0
2	d1			d1	a1	d0	c1	a0
1	c0	b1		b0	c0		a1	c0
0	a1	d0	c1	a0	d0	b1	d1	b0
	0	1	2	3	4	5	6	7

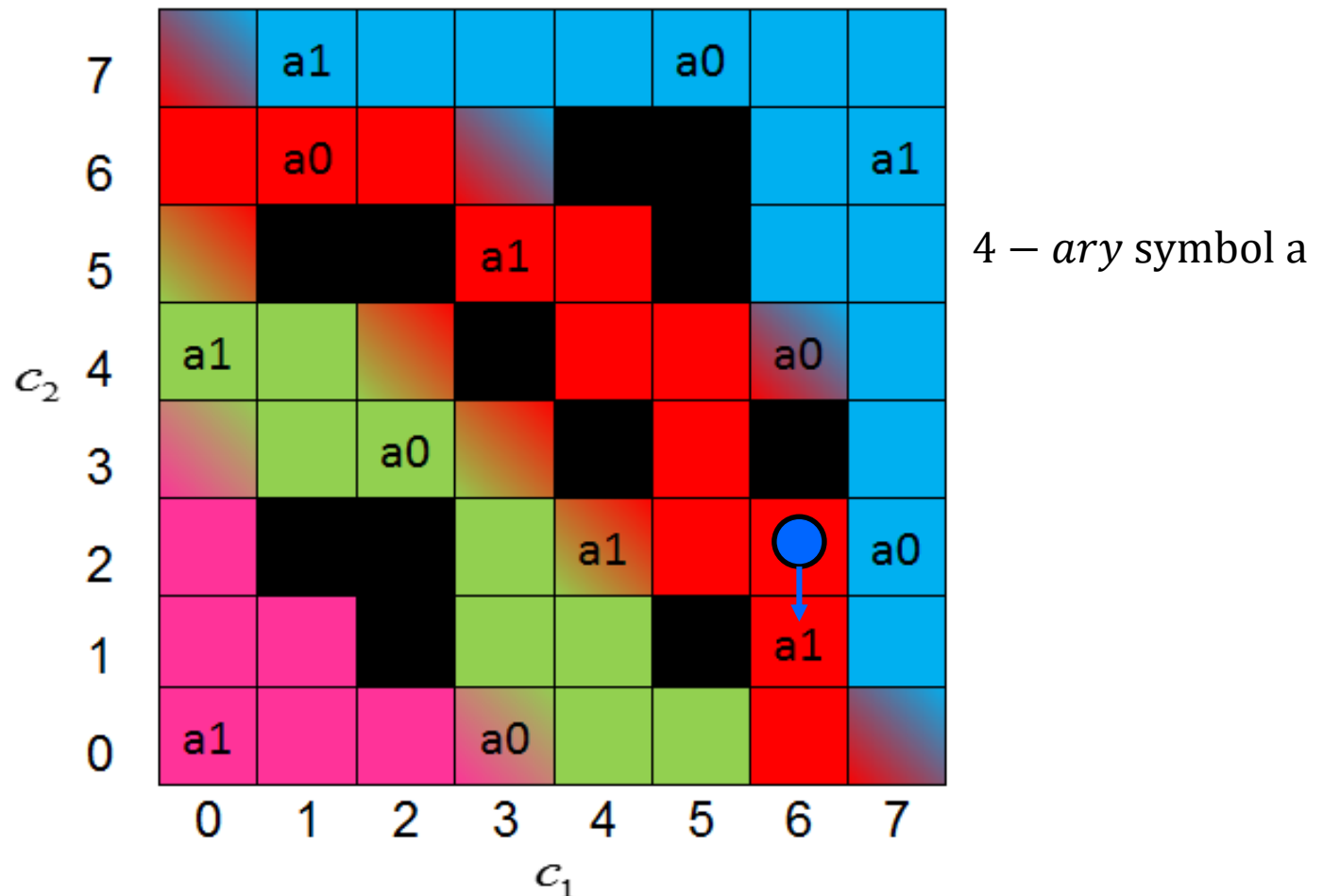
c_2

c_1

EC-WOM construction for mag-1 errors



- Decoder knows correct 4 – ary symbol and write number
- Looks for closet 4 – ary symbol within same write region



Comparison of Constructions



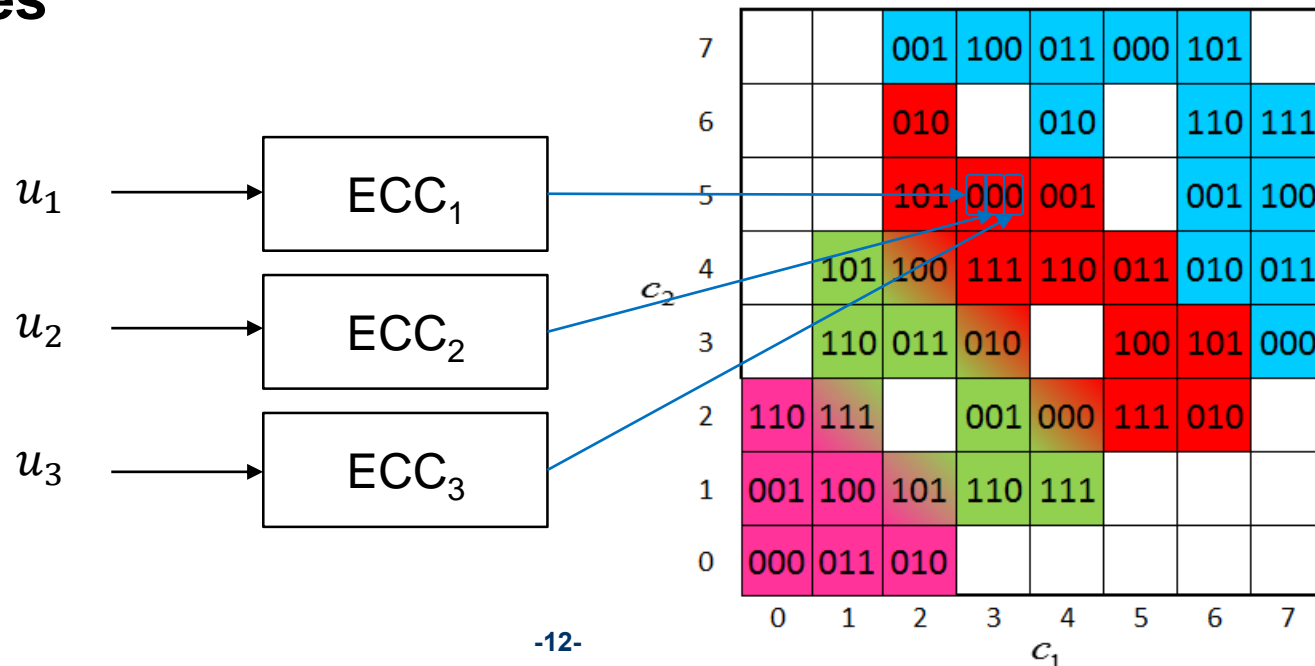
■ Closed form $\forall \tau, N$ (approximate)

Construction	EC-Rate	Error Type	Alphabets
1	$1 - \frac{\frac{7}{8}(2\tau - 1)\log_8(N)}{N}$	Any	8
2	$1 - \frac{\left(\frac{3}{4}(\tau - 1) + \frac{1}{2}(2\tau - 1)\right)\log_8 N}{N}$	mag-1	2,4
3'	$1 - \frac{\frac{1}{4}(6\tau_1 + 10\tau_2 - 5)\log_8 N}{N}$	τ_1 single mag-1, τ_2 general	4,2
4	$1 - \frac{\frac{3}{2}(3\tau - 2)\log_8 N + 2\tau - 1}{N}$	amag-1	2,2

EC-WOM codes for random errors



- **Objective** – improve and analyze WOM error correction over realistic memory channels
- **Multi level coding** [Wachsmann, Fischer, Huber '99]
- **High order symbols are mapped to multiple bits**
- **Main idea: propagate information from low to high bits**
- **Multi-stage decoding, WOM converts many errors to erasures**



Erasures



Erasures: write number $k = 3$, bit number 1 $i = 1$ (LSB)

7			001	100	011	000	101	
6			010		010		110	111
5			101	000	001		001	100
4		101	100	111	110	011	010	011
3		110	011	010		100	101	000
2	110	111		001	000	111	010	
1	001	100	101	110	111			
0	000	011	010					
	0	1	2	3	4	5	6	7

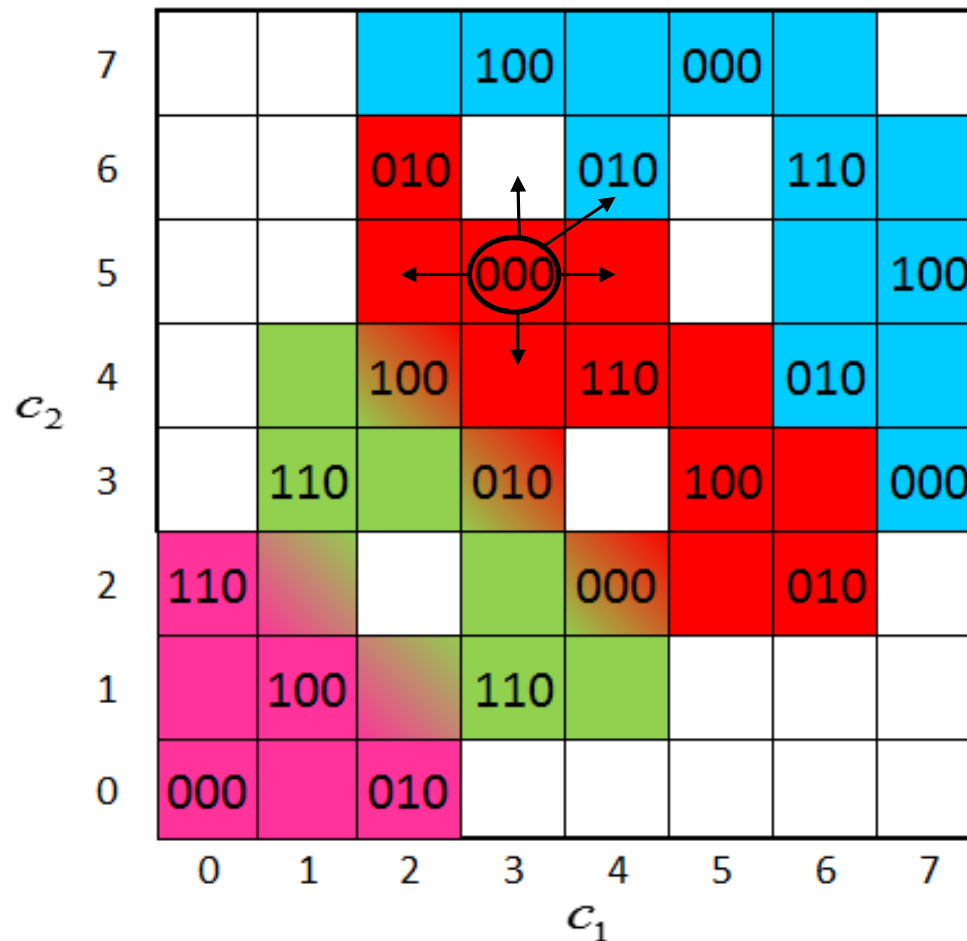
c_2

c_1

Erasures



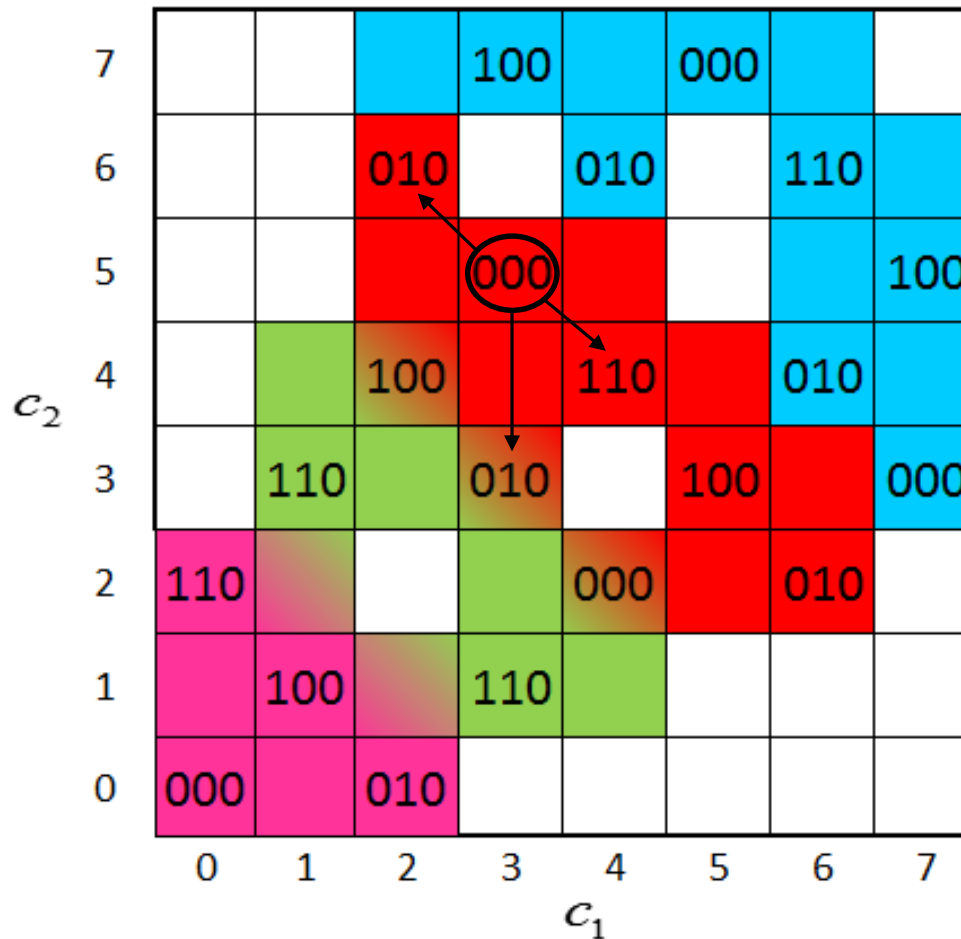
Erasures: write number $k = 3$, bit number $i = 2$, first bit $b_1 = 0$



Errors

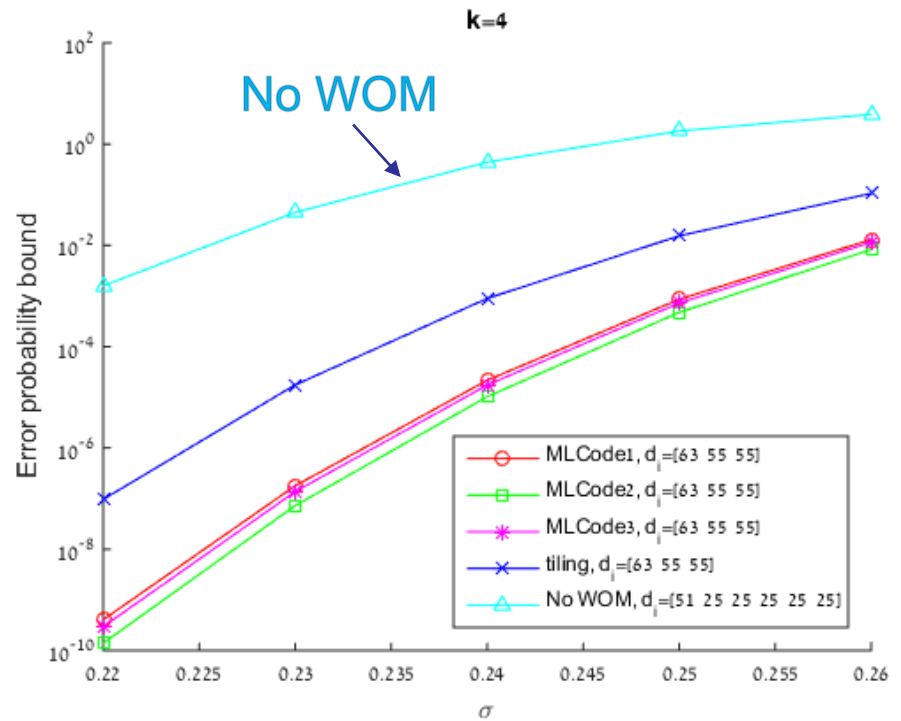
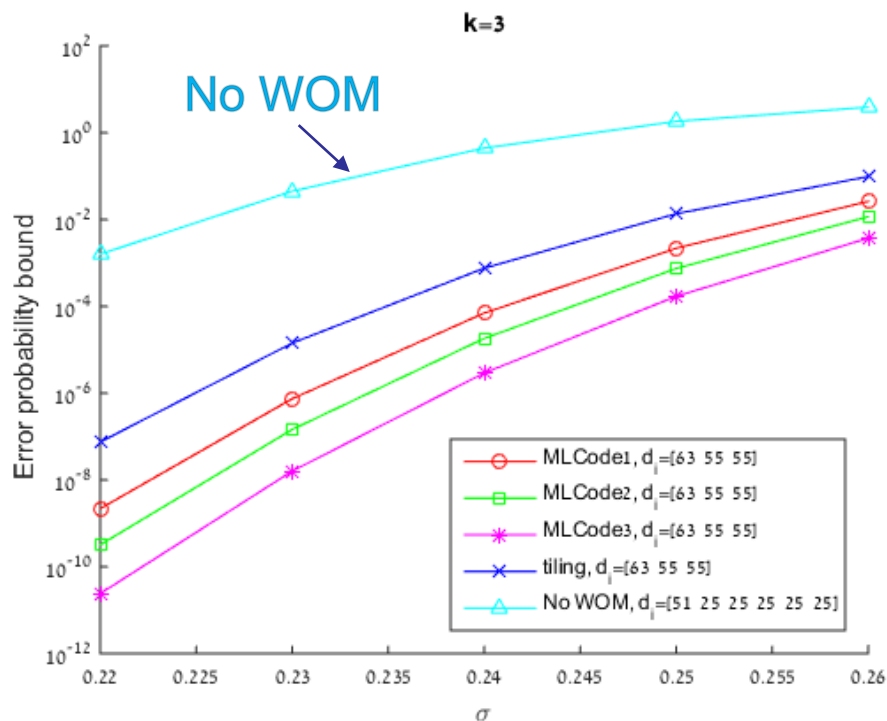


Errors: write number $k = 3$, bit number $i = 2$, first bit $b_1 = 0$



Error probability bound – AWGN example

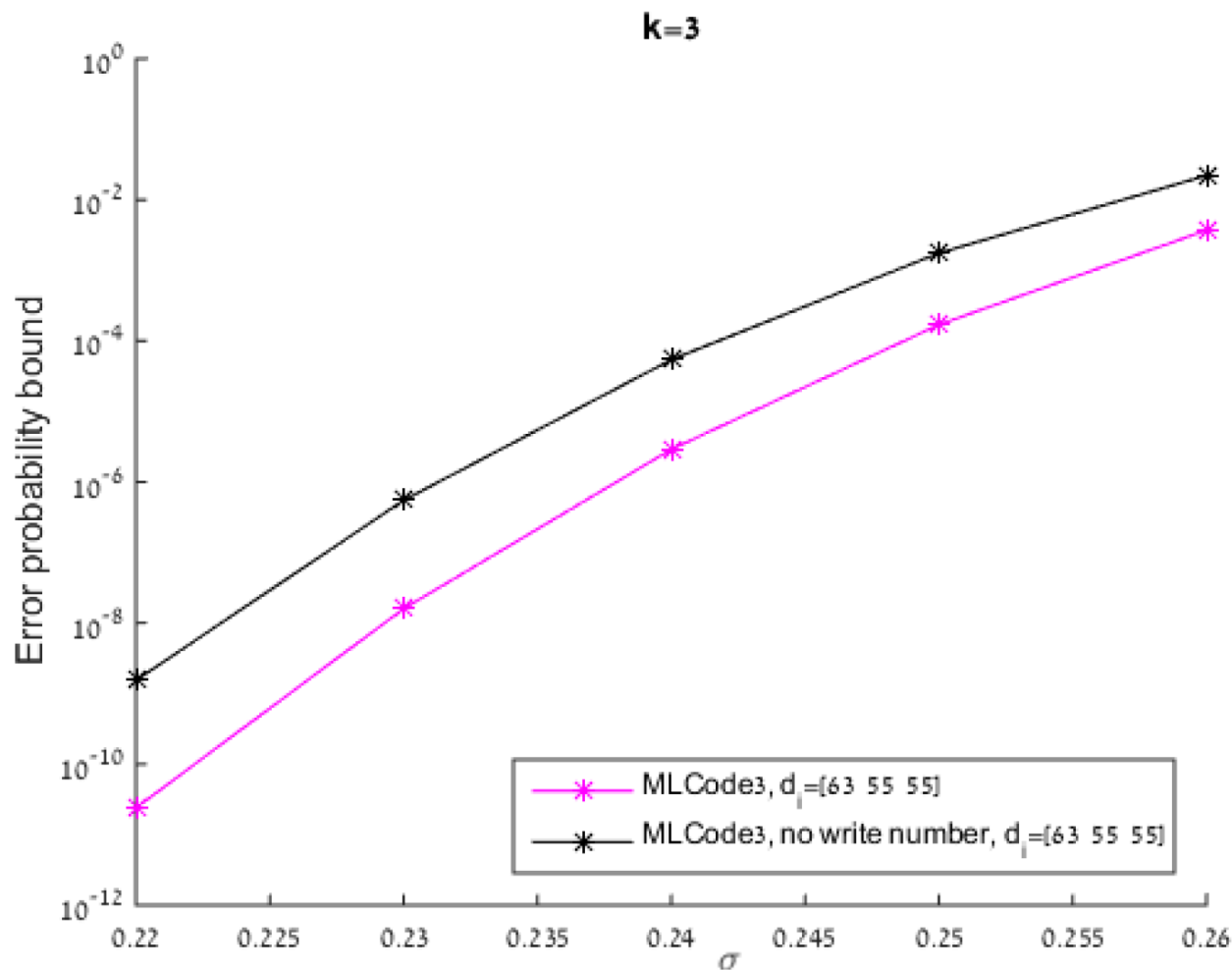
- AWGN analysis, standard deviation σ
- Optimized redundancy distribution
- Code performance changes with write number k
- Propagation Index can predict WOM performance



Knowing the write number helps



- Comparison between decoders with/without write-number information



Multi level coding soft decoding



- Channel output has additional reliability information
- Use soft BCH decoders: GMD [Forney '73]
- Reliability 1 – Certain bit
- Reliability 0 – Erasure
- **Soft information with hard memory reads**

1	1	1
1	0	1
1	1	1

Low
reliability

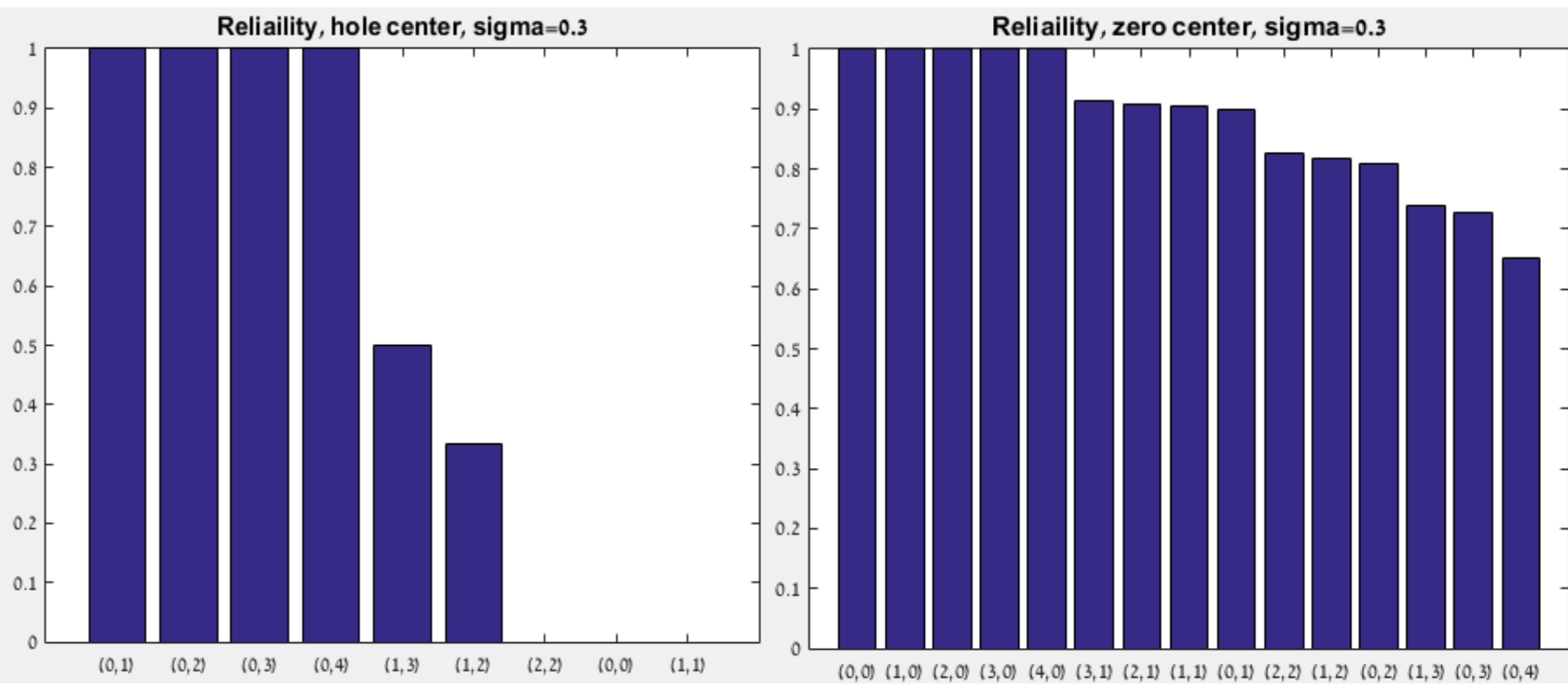
0	0	0
0	0	0
0	0	0

High
reliability

Multi level coding soft decoding - reliability



- Problem: order reliabilities of each physical state
- AWGN: *non* mag-1 errors are unlikely
- (m, n) : m adjacent zeroes, n adjacent ones
- Lemma: order of reliabilities does not change for $\sigma < 0.341$





- **A. Solomon and Y. Cassuto, “Error-Correcting WOM Codes: Concatenation and Joint Design”, IEEE Transactions on Information Theory, 09/2019.**
- **A. Solomon and Y. Cassuto, “Error-Correcting WOM Constructions through Concatenation and Joint Design”, IEEE International Symposium on Information Theory, ISIT 2018.**

- WOM codes can help flash based memories
- Joint design of WOM and ECC codes improves reliability
 - Guaranteed error models
 - Random error models
- Also improves decoding complexity
- Can obtain soft information with hard memory reads

*Thank
you*

